

W13D1
18/1/15

Prob. Methods in Comb.

Thm. $P \sim \text{Pois}(\mu)$, then (i) $P(P \leq (1-\epsilon)\mu) \leq e^{-\epsilon^2 \mu/2}$, $P(P \geq (1+\epsilon)\mu) \leq [e^{\epsilon(1-\epsilon)}]^\mu$.

Recall We construct $R \subseteq \Omega$ at random by $P(x \in R) = p_x$ indep. of other x . $B_i = \{x_i \in R\}$. $\mu = \sum_{i \in I} P(B_i) = \sum_{i=1}^n \prod_{x \in A_i} p_x$.

$\Delta = \sum_{\substack{(i,j) \in I^2 \\ i \sim j}} P(B_i \cap B_j)$ where $i \sim j \Leftrightarrow i \neq j$ and $A_i \cap A_j \neq \emptyset$.

$J \subseteq I$ is a maximal disjoint family $\Leftrightarrow (\bigcap_{i \in J} B_i \text{ holds and } \forall i \notin J, B_i \rightarrow \text{some } j \in J)$.

For $s \in \mathbb{N}$, $\mu_s = \min \{ \sum_{i \in J} P(B_i) \mid J \subseteq I \text{ of size } s \}$. $\nu = \max_{j \in I} \sum_{i \sim j} P(B_i)$. Then observe $\mu_s \geq \mu - s \cdot \nu$.

Lemma $P(\exists \text{ max. disj. family } J \text{ of size } s) \leq \frac{\nu^s}{s!} e^{-\mu} e^{\Delta/2}$.

pf. Given $J \subseteq I$ of size s , $i \not\sim j$ for all $i, j \in J$, $P(J \text{ is max. disj. family}) = P(\bigcap_{i \in J} B_i \cap \bigcap_{i \notin J} \bar{B}_i) = \prod_{i \in J} P(B_i) \cdot P(\bigcap_{i \notin J} \bar{B}_i) \leq \prod_{i \in J} P(B_i) e^{-\mu(J) \Delta(J)/2} \leq$

$\prod_{i \in J} P(B_i) e^{-\mu_s + \Delta/2}$. So, $P(\exists \text{ such family}) \leq \sum_{\substack{J \subseteq I \\ |J|=s}} \prod_{i \in J} P(B_i) e^{-\mu_s + \Delta/2} \leq \frac{\nu^s}{s!} e^{-\mu_s + \Delta/2}$.

Note Previously, we could prove $P(\exists \text{ max. disj. fam. of size } \geq 3\mu) \leq$

$P(\exists \text{ disj. family of size } 3\mu) \leq (\frac{e}{3})^\mu$. If $\Delta = o(1)$ and $\mu_{3\mu} = \mu - o(1)$

then for $s \leq 3\mu$ we get $P(\exists \text{ max. disj. fam. of size } s) = (1 + o(1)) \frac{\mu^s}{s!} e^{\Delta/2} = (1 + o(1)) P(\text{Pois}(\mu) = s)$.

Every vertex of $G(n, p)$ is in a Triangle (Rev.)

Let $N(x) = \# \text{ triangles which contain } x$. $\mathbb{E} N(x) = \binom{n-1}{2} p^3 = \mu$.

Thm Let p be such that $\mu \gg \log n$. Then a.a.s., $N(x) = (1 + o(1))\mu$ for every x . More precisely, $P(|N(x) - \mu| \geq \epsilon \mu) = o(n^{-1})$ for every $x \in V(G)$, $\epsilon > 0$ (and the result follows from the union bound).

pf. For simplicity assume $\mu = n^{o(1)}$ (eg. $p = n^{-2/3 + o(1)}$). $\Delta = O(n p^5) = o(1)$, $\nu = O(n p^3)$, hence $\mu \nu = o(1)$, $\mu_s = \mu - o(1)$. Define $\tilde{N}(x) = \text{max. \# edge-disjoint } \Delta \text{ s containing } x$. Clearly $N(x) \geq \tilde{N}(x)$ and $P(N(x) \leq (1-\epsilon)\mu) \leq P(\tilde{N}(x) \leq (1-\epsilon)\mu) = (1 + o(1)) P(\text{Pois}(\mu) \leq (1-\epsilon)\mu) \leq (1 + o(1)) e^{-\epsilon^2 \mu/2} \leq o(n^{-2})$ for large enough n . For the upper tail,

$$P(\tilde{N}(x) \geq (1-\epsilon)\mu) = \sum_{s=(1-\epsilon)\mu}^{\mu} P(\tilde{N}(x)=s) + P(\tilde{N}(x) > 3\mu) \leq \frac{1}{n^2} + \left(\frac{e}{3}\right)^\mu \leq \frac{1}{n^2} + \frac{1}{n^2} \text{ for large } n$$

enough n , since $P(\text{Pois}(\mu) \geq (1-\epsilon)\mu) \leq (1-\delta_\epsilon)^\mu \xrightarrow{n \rightarrow \infty} 0$.

$$P(\exists \text{ triangle}) \leq n^3 p^3 = o(n^{-1}), P(\exists \text{ cycle of length } \leq 4) \leq n^4 p^4 = o(n^{-1})$$

Consider an aux. graph T with vertices \leftrightarrow triangles containing x , edge \leftrightarrow common edge. We proved w. prob $1 - o(n^{-1})$: $\alpha(T) \leq (1+\epsilon)\mu$, $\Delta(T) \leq 3$, induced max. matching ≤ 3 . We claim $\alpha(T) \geq V(T) - 27$

so $V(T) \leq \alpha(T) + 27 \leq (1+2\epsilon)\mu$. Why is that? While \exists edge, delete

both endpoints and their neighbours (≤ 6), after 3 iterations we can't have other edges (otherwise, we'd have an

induced matching of size 4).

X r.v. $\in R$ finite set.

Self-Information is a measure of the information content associated with the outcome of $X \in S$, $I(S)$ - info. content of $x \in S$.

(i) $I(S)$ depends only on $P(x \in S)$, (ii) If $x \in S_1$ and $x \in S_2$ are indep., $I(S_1 \cap S_2) = I(S_1) + I(S_2)$, $I(S)$ nonzero.

Only choice is $\log_b \frac{1}{P(x \in S)}$ for some $b > 1$. $b=2$ "bits",

$b=e$ "nats", $b=10$ "bans/dits/hartleys".

Entropy is the "average" information content of an outcome:

$$H(X) = \mathbb{E} I(x) = \sum_{x \in R} P(X=x) \log_2 \frac{1}{P(X=x)}$$

Conditional Entropy If $X, Y: \Omega \rightarrow R, S$ r.v. then $(X, Y): \Omega \rightarrow R \times S$.

$$\text{Then } H(X|Y) \stackrel{\text{def}}{=} H(X, Y) - H(Y)$$

Lemma X, Y, Z r.v. - (i) $0 \leq H(X) \leq \log_2 |\text{range}(X)|$

$$(ii) H(X, Y) \geq H(X)$$

$$(iii) H(X, Y) \leq H(X) + H(Y)$$

$$(iv) H(X|Y, Z) \leq H(X|Y)$$

pf. (i) $z \mapsto \log_2 z$ is concave, hence $\sum_{x \in R} P(X=x) \log_2 \frac{1}{P(X=x)} \leq$

$$\log_2 \frac{1}{\sum_{x \in R} P(X=x)} = \log_2 \frac{1}{1} = \log_2 |\text{range}(X)|. \text{ Sharp iff } P(X=x) = \frac{1}{|\text{range}(X)|} (\forall x \in R).$$

$$(ii) H(X, Y) = \sum_{x \in R} \sum_{y \in S} P(X=x, Y=y) \log_2 \frac{1}{P(X=x, Y=y)} \leq \sum_{x \in R} \sum_{y \in S} P(X=x, Y=y) \frac{1}{P(X=x)} = \sum_{x \in R} P(X=x) \log_2 \frac{1}{P(X=x)} \stackrel{\text{Jensen}}{=} H(X)$$

$$(iii) H(X) + H(Y) - H(X, Y) = \sum_{x \in R} \sum_{y \in S} P(X=x, Y=y) \log_2 \frac{P(X=x)P(Y=y)}{P(X=x, Y=y)} = \sum_{x \in R} \sum_{y \in S} P(X=x)P(Y=y) f(z_{x,y}) \stackrel{\text{Jensen}}{\geq} f(\sum_{x \in R} \sum_{y \in S} P(X=x, Y=y)) = f(1) = 0$$

$f(z) = z \log_2 z$ convex

(iv) will be shown later.

Prop. Let $x = (x_1, \dots, x_n): \Omega \rightarrow \mathbb{R}_1 \times \dots \times \mathbb{R}_n$. By applying (iii) repeatedly, $H(x) \leq \sum_{i=1}^n H(x_i)$.

cor. Let \mathcal{F} be a family of subsets of $[n]$, $p_i =$ fraction of sets in \mathcal{F} which contain i . Then $|\mathcal{F}| \leq 2^{\sum_{i=1}^n H(p_i)}$ where

$$H(p) = H(\text{Bin}(p)) = p \log_2 \frac{1}{p} + (1-p) \log_2 \frac{1}{1-p}$$

pf. Choose $X \in \mathcal{F}$ u.a.r. $- H(x) = \log_2 |\mathcal{F}| \leq H(x_1) + \dots + H(x_n) = \sum_{i=1}^n H(p_i)$

cor. $\forall n \geq 1, \sum_{k=0}^n \binom{n}{k} \leq 2^{nH(1/2)}$ sharp, up to poly(n) factor. $\binom{n}{\lfloor n/2 \rfloor} = \frac{2^n}{\sqrt{2\pi n} \cdot (1/2)^n}$

pf. Take \mathcal{F} to be the family of subsets of size $\leq \lfloor cn \rfloor$ of $[n]$. $p_i \leq c$ and H is increasing on $[0, 1/2]$.

cor Let $X_1, \dots, X_n \sim U[-1, 1]$ i.i.d. Then $\forall c > 0, P(|X| \geq cn) \leq 2^{-c^2 n}$ where $X = \sum_{i=1}^n X_i$.

Lemma (Shearer '86) Let $X = (x_1, \dots, x_n)$ r.v. $\mathcal{G} \subseteq \mathcal{P}([n])$ such that any $i \in [n]$ belongs to at least k $\mathcal{G} \in \mathcal{G}$. Then $H(X) \leq \frac{1}{k} \sum_{\mathcal{G} \in \mathcal{G}} H(X(\mathcal{G}))$ where $X(\mathcal{G}) = (x_i)_{i \in \mathcal{G}}$.

We'll probably only show this next week due to time constraints.

Cor $\mathcal{F} \subseteq S_1 \times \dots \times S_n$, $\mathcal{G} = \{G_1, \dots, G_n\}$ such that each $i \in [n]$ belongs to at least k $\mathcal{G} \in \mathcal{G}$. Let F_i be the proj. of \mathcal{F} onto G_i .

$$\text{Then } |\mathcal{F}|^k \leq \prod_{i \in [n]} |F_i|$$

pf. $X \sim \text{Uni}(\mathcal{F})$, then by Shearer $\log_2 |\mathcal{F}| = H(x) \leq \frac{1}{k} \sum_{i \in [n]} \log_2 |F_i|$

Cor Let $B \subseteq \mathbb{R}^n$ measurable, $\forall i: B_i = \Pi_i(B)$ (The projection to $x_i = 0$). Then

$$\text{Loomis-Whitney '49} \quad \text{Vol}_n(B) \leq \prod_{i=1}^n \text{Vol}_{n-1}(B_i)$$

pf Approximate B using ϵ -cubes, $\mathcal{F} =$ set of cubes, use prev. cor.

cor. Let \mathcal{F} be a family of graphs on $V = [n]$ s.t. $\forall G_1, G_2 \in \mathcal{F}$,

$$K_3 \subseteq G_1 \cap G_2. \text{ Then } |\mathcal{F}| \leq \frac{c}{4} 2^{\binom{n}{2}}$$

$$sm = k \binom{n}{2}$$

pf. Let $\mathcal{G} = \{A \subseteq [n], |A| = \lfloor n/2 \rfloor\}$. $\forall \mathcal{G} \in \mathcal{G}$, $\Pi_{\mathcal{G}}(\mathcal{F})$ intersecting, hence of size $\leq 2^{e(\mathcal{G})-1}$.

sketch Each edge is covered $\leq k$ times by \mathcal{G} . By cor. above, $|\mathcal{F}| \leq$

$$\prod_{\mathcal{G} \in \mathcal{G}} 2^{e(\mathcal{G})-1} = \frac{1}{2} \prod_{\mathcal{G} \in \mathcal{G}} 2^{e(\mathcal{G})} = \frac{1}{2} 2^{sm} = \frac{1}{2} 2^{k \binom{n}{2}} = \frac{1}{2} 2^{\binom{n}{2}}$$