

Prob. Methods in Comb.

Talagrand's inequality -  $\Omega_i$  prob. space,  $\Omega = \prod_{i=1}^n \Omega_i$  with product measure.

Thm (Talagrand) For any  $A \subseteq \Omega$ ,  $x \in \Omega$  random,  $t \geq 0$ ,  
 $P(\sum_{i=1}^n x_i \in A) P(\rho(x, A) \geq t) \leq e^{-t^2/4}$  where  $\rho(x, A) = \sup_{\|x\|_2=1} \min_{y \in A} \sum_{i: x_i \neq y_i} \alpha_i$ .

Example longest increasing subsequence of a random permutation.  
 In HW - LIS( $\sigma$ ) =  $\Theta(\sqrt{n})$  a. a. s. when  $\sigma \in S_n$  u. a. r.

We show a. a. s. LIS( $\sigma$ ) belongs to an interval of length  $w(n) \cdot n^{1/4}$  for any  $w(n) \rightarrow \infty$ .

Fix  $f: \mathbb{N} \rightarrow \mathbb{N}$ .  $h$  is  $f$ -certifiable if whenever  $h(x) \geq s$   
 $\exists I_x \subseteq [n]$ ,  $|I_x| \leq f(s)$  and  $\forall y \in \Omega$  s.t.  $x|_{I_x} = y|_{I_x} \Rightarrow h(y) \geq s$ .

Ex  $\Omega = \{0, 1\}^{\binom{n}{2}}$ ,  $h = \# \Delta$  in  $G$ ,  $h$  is  $f$ -cert. for  $f(s) = 3s$ .

Ex If  $f(s) = n$  any  $h$  is  $f$ -cert.

Thm Let  $h: \Omega \rightarrow \mathbb{R}$  be  $f$ -cert. Lipschitz function.

Let  $x \in \Omega$  random, look at  $h(x)$ . Then  $\forall b, t$  -  
 $P(h(x) \geq b) P(h(x) \leq b - t \sqrt{f(b)}) \leq e^{-t^2/4}$ .

pf  $A = \{y \in \Omega \mid h(y) < b - t \sqrt{f(b)}\}$ . Suppose  $h(x) \geq b$  - it's enough to show  $\rho(x, A) \geq t$  and apply Talagrand. If  $h(x) \geq b \exists I_x \dots$   
 Set  $\alpha_i = \begin{cases} |I_x|^{-1/2} & i \in I_x \\ 0 & \text{o/w} \end{cases}$  so that  $\|\alpha\|_2 = 1$ . So  $d_\alpha(x, A) < t$  and  $\exists y \in A$  with  $d_\alpha(x, y) < t$ . Define  $x'_i = \begin{cases} x_i & i \in I_x \\ y_i & \text{o/w} \end{cases}$ .  
 Then  $h(x) \geq b$  by def of  $I_x$  -  $b d_\alpha(y, x) = \frac{1}{|I_x|^{1/2}} \# \{x_i \neq y_i\}$ . So  $d_\alpha(x, y) = \# \{i \in I_x \mid x_i \neq y_i\} < t \cdot |I_x|^{1/2} \leq t \sqrt{f(b)}$  but then  $h(y) \geq h(x) - d_\alpha(x, y) \geq b - t \sqrt{f(b)}$  - but  $y \in A$  so  $h(y) < b - t \sqrt{f(b)}$  in contradiction.

For our setting, note that replacing random  $\sigma \in S_n$  with  $x \in \{0, 1\}^{\binom{n}{2}}$  doesn't change the dist. of LIS( $x$ ). It's Lipschitz and  $(s \rightarrow 3s)$  certifiable. Let  $w = b = \text{median of LIS}(x) = \Theta(\sqrt{n})$ . Then

then  $\mathbb{P}(\text{LIS}(x) \geq m) \mathbb{P}(\text{LIS}(x) \leq m - t\sqrt{m}) \leq e^{-t^2/4}$ . For the upper bound  $m = b - t\sqrt{b}$  we get  $\mathbb{P}(\text{LIS}(x) - m \leq t\sqrt{m}) \leq 4e^{-t^2/4}$ . Bounded differences only gives  $2e^{-2t^2} \leq \mathbb{P}(\text{LIS} - \mu \leq t\sqrt{n})$ .

Sketch of pf. of Talagrand:

$$\mathbb{P}(\rho(x, A) \geq t) = \mathbb{P}(e^{\frac{\rho^2(x, A)}{4}} \geq e^{t^2/4}) \stackrel{\text{Markov}}{\leq} e^{-t^2/4} \mathbb{E}[e^{\frac{\rho^2(x, A)}{4}}]$$

(\*) can be proved inductively on  $n$ .  $(*) \rightarrow \sum_{x \in A} 1$

### Today: Vanson's inequality

Motivation We want to estimate  $\mathbb{P}(\bigcap_{i \in I} \bar{B}_i)$ . Vanson gives an upper bound.

Note - If  $B_i$  were mutually ind. was just  $\prod_{i \in I} \mathbb{P}(\bar{B}_i) \approx e^{-\mathbb{E}[\#B_i \text{ that hold}]}$  for  $\mathbb{P}(B) \ll 1$ . Jensen - under certain assumption this is close to the truth.

Setting  $\Omega$  fin. For each  $x \in \Omega$  fix  $p_x \in [0, 1]$  and let  $R \subseteq \Omega$  random, where  $\forall x \mathbb{P}(x \in R) = p_x$  indep. of  $x' \neq x \in \Omega$ .

Fix  $A_i \subseteq \Omega$  and let  $B_i = \{A_i \subseteq R\}$ . For  $i, j \in I$  write  $i \sim j$  iff  $i \neq j$  but  $A_i \cap A_j \neq \emptyset$ . Define  $M = \prod_{i \in I} (1 - \mathbb{P}(B_i))$ ,

$$\mu = \mathbb{E}[\#B_i \text{ that hold}] = \sum_{i \in I} \mathbb{P}(B_i), \Delta = \mathbb{E}[(i, j) \in I^2 \mid i \sim j, B_i, B_j \text{ hold}] =$$

$$\sum_{i \sim j} \mathbb{P}(B_i \cap B_j) = \sum_{i \sim j} \prod_{x \in A_i \cap A_j} p_x$$

Vanson In the above notation,  $M \leq \mathbb{P}(\bigcap_{i \in I} \bar{B}_i) \leq e^{-\mu + \Delta/2}$ . If moreover,  $\forall i \mathbb{P}(B_i) \leq \epsilon$ ,  $\mathbb{P}(\bigcap_{i \in I} \bar{B}_i) \leq M \cdot e^{\frac{\Delta}{2\epsilon}}$ .

comment If  $\epsilon \Delta = o(1)$  then  $\mathbb{P}(\bigcap_{i \in I} \bar{B}_i) = (1 - o(1)) e^{-\mu} = (1 - o(1)) M$ .

If  $\Delta \geq 2\mu$ , (\*) is useless but Thm (Extended Vanson) If

$$\Delta \geq \mu, \mathbb{P}(\bigcap_{i \in I} \bar{B}_i) \leq e^{-\mu/2\Delta}$$

pf. Given  $J \subseteq I$  define  $\mu(J), \Delta(J)$  like we did with  $I$ . Let  $q \in [0, 1]$  and  $J$  be a random subset of  $I$ , keeping each element with prob.  $q$ .  $\mathbb{E}[\mu(J) - \frac{\Delta(J)}{2}] = q\mu - \frac{1}{2}q^2\Delta = -\frac{\mu^2}{2\Delta}$ . Hence  $\exists J$  with  $-\mu(J) - \Delta(J)/2 \leq -\frac{\mu^2}{2\Delta}$  so we get the bound by applying (regular) Vanson.

Vanson Ex. Vanson

pf. Assume  $I = \{1, \dots, m\}$ .  $P(\bigcap_{i=1}^m \overline{B}_i) = \prod_{i=1}^m P(\overline{B}_i | \bigcap_{j=1}^{i-1} \overline{B}_j)$ . We claim  $P(B_i | \bigcap_{j=1}^{i-1} \overline{B}_j) \leq P(B_i)$ . Claim  $\Rightarrow P(\bigcap_{i=1}^m \overline{B}_i) \geq \prod_{i=1}^m (1 - P(B_i) + \sum_{j=1}^{i-1} P(B_i | \bigcap_{j=1}^{j-1} \overline{B}_j))$  which is  $\leq \prod_{i=1}^m e^{-P(B_i) + \sum_{j=1}^{i-1} P(B_i | \bigcap_{j=1}^{j-1} \overline{B}_j)} = e^{-\mu + \Delta/2}$  but also (if  $\forall i: P(B_i) \leq \epsilon$ )  $\leq \prod_{i=1}^m [1 - P(B_i)] \left[ 1 + \frac{\sum_{j=1}^{i-1} P(B_i | \bigcap_{j=1}^{j-1} \overline{B}_j)}{1 - \epsilon} \right] \leq n \cdot e^{\frac{\Delta}{2(1-\epsilon)}}$ .

pf. claim  $D_i = \bigcap_{j=1}^d \overline{B}_j, U_i = \bigcap_{j=d+1}^m \overline{B}_j, P(B_j | D_i \cap U_i) \geq P(B_j | U_i) = P(B_j)$  (we assume  $\{j \sim i\} = \{1, \dots, d\}$ )  $= P(B_j | U_i) P(D_i | B_j \cap U_i)$ .  $P(D_i | B_j \cap U_i) = 1 - P(\bigcup_{j=1}^d B_j | B_j \cap U_i) \geq 1 - \sum_{j=1}^d P(B_j | B_j \cap U_i) \geq 1 - \sum_{j=1}^d P(B_j | B_j)$ . It's enough to prove (\*) term by term.  $P(B_j | B_j \cap U_i) = P^*(B_j | U_i) \leq P^*(B_j) = P(B_j | B_j)$ .

Ex. clique num of  $G(n, 1/2)$ . Let  $f(k) = \binom{n}{k} 2^{-\binom{k}{2}}$ . Def.  $k_0 = \min\{k | f(k) < 1\}$ .  $f(k_0 + 1) \approx n^{-1+o(1)}$  so a. e. s.  $w(G) < k_0 + 1$ .  $k_1 = k_0 - 4, f(k_1) = n^{4+o(1)}$ .

and we proved  $P(w(G) < k_1) \leq e^{-c \frac{n^2}{\log^4 n}}$ . We can now improve to  $\log^4 n$ .  $I = \binom{[n]}{k}$ . Given  $S \in I$  define  $B_S = \{G[S] \text{ clique}\}$ . Then  $P(w(G) \geq k_1) = P(\bigcup_{S \in I} B_S)$ .  $\mu = f(k_1) \geq n^{3+o(1)}$ ,  $\Delta = \sum_{S \in I} \sum_{T \in I} P(B_S \cap B_T) = \sum_{s=2}^k \binom{n}{s} \binom{n-k}{k-s} \binom{k-s}{2} 2^{-\binom{k-s}{2}} = \sum_{s=2}^k \binom{n}{s} \binom{k-s}{2} 2^{-\binom{k-s}{2}}$ .  $(1+o(1)) \frac{f^2(k_1) k_1^4}{n^2} = (1+o(1)) \frac{\mu^2 k_1^4}{n^2} \gg \mu$ . By EYI  $P(w(G) \geq k_1) \leq e^{-\frac{\mu^2}{2\Delta}} \leq e^{-c \frac{n^2}{\log^4 n}} = e^{-c' \frac{n^2}{\log^4 n}}$ .

Ex. Triangles in  $G(n, p)$  - what is  $P(G \text{ is } \Delta\text{-free})$ ? For large  $p$  this is  $\approx P(G \text{ bipartite})$ . But what if  $p = \frac{c}{n}$ ?  $X$  - num of  $\Delta$  in  $G$ .  $E[X] = (1+o(1)) \frac{c^3}{6}$  [In fact,  $X \approx \text{Pois}(\frac{c^3}{6})$ ].

claim  $P(X=0) \approx e^{-c^3/6}$ . pf. Yanson with  $I = \binom{[n]}{3}, B_S = \{G[S] \approx K_3\}$ . Then  $\mu = \binom{n}{3} p^3 = (1+o(1)) \frac{c^3}{6}$ .  $\Delta = \binom{n}{3} (n-3) \cdot 3 \cdot p^5 = \Theta(n^4 p^5) = \Theta(\frac{1}{n}) = o(1)$ . By Yanson FKG  $(1+o(1)) e^{-\frac{c^3}{6}} = (1-p^3)^{\binom{n}{3}} \leq P(X=0) \leq e^{-\mu + o(1)} = (1+o(1)) e^{-\frac{c^3}{6}}$ .