

I don't know

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! module

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Abstract Data Type

length

Retrieve

some loop

when in

(well F!)

↑

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ADT - a
- b

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a) Remove Evens (L)

for i = 0 to Length(L) - 1

if is even (retrieve(L, i)) then delete(L, i)

b) Remove evens (L):

while (L.first ≠ null) and in even L.first.item)

L.first ← L.first.next

Decrease L.length

~~A ← L.first~~

A ← L.first

return

If A = Null then return

return next → 30

(return), (return)

(return) (return) no

while A.next != Null

if is even(A.next.item) then A.next ← A.next.next

else A ← A.next

decrease
length

$$f(n) = O(g(n)) \iff \exists c, n_0 \forall n > n_0 f(n) \leq c \cdot g(n)$$

$$f(n) = \Omega(g(n)) \iff \exists c, n_0 \forall n > n_0 f(n) \geq c \cdot g(n)$$

$$f(n) = \Theta(g(n)) \iff f = O(g(n)) \wedge f = \Omega(g(n))$$

$$f(n) = o(g(n)) \iff \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \rightarrow 0$$

possible

$$\forall \epsilon, \exists n_0 \forall n > n_0$$

$$c \cdot f(n) \leq g(n)$$

$$f(n) = \omega(g(n)) \iff \forall M > 0 \exists n_0 \forall n > n_0 \frac{f(n)}{g(n)} > M$$

$$n^3 = o(n^4)$$

$$2^n = \Omega(3^n)$$

$$n = \omega(\log n)$$

$$10e^n = \Theta(e^n)$$

~~f(n) = 0~~

$$\limsup_{n \rightarrow \infty} \frac{f(n)}{g(n)} < P \quad f = O(g)$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \quad f = o(g)$$

$$\liminf_{n \rightarrow \infty} \frac{f(n)}{g(n)} > 0 \quad f = \Omega(g)$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty \quad f = \omega(g)$$

$$O(b^h) \quad b \geq 2$$

$$O(2^h)$$

$$O(h^\alpha) \quad \alpha \geq 2$$

$$O(h^2)$$

$$O(h \log h)$$

$$O(h)$$

$$O(\sqrt{h})$$

$$O(\log h)$$

$$O(1)$$

↖ ↗ ↘ ↙

$$\Theta(h^{0.99})$$

$$\Theta(h \log h)$$

$$\Theta\left(\frac{h}{\log h}\right)$$

$$\Theta(h^{0.99})$$

$$\leq \Theta\left(\frac{h}{\log h}\right)$$

$$\leq \Theta(h \log h)$$

$$\frac{h}{\log h}$$

$$= \frac{h^{0.01}}{\log h}$$

Ⓛ

$$= \frac{0.01}{h^{0.99}} \cdot \frac{1}{\frac{1}{h}}$$

$$= 0.01 \cdot h^{0.01}$$

$$\rightarrow \infty$$

$$f(h) = \Theta(g(h))$$

$$\log(f(h)) \stackrel{?}{=} \Theta(\log(g(h)))$$

$$f(h) = c \cdot g(h)$$

$$\log f(h) = \log f(h) + \log c$$

$$2^{f(h)} \stackrel{?}{=} \Theta(2^{g(h)})$$

$$g(h) = 2^h \quad f(h) = 2^h$$

↖ ↗ ↘ ↙

$$\lim \frac{f}{g} = \lim \frac{2^{2^h}}{2^h} = 2^h \neq \text{const}$$